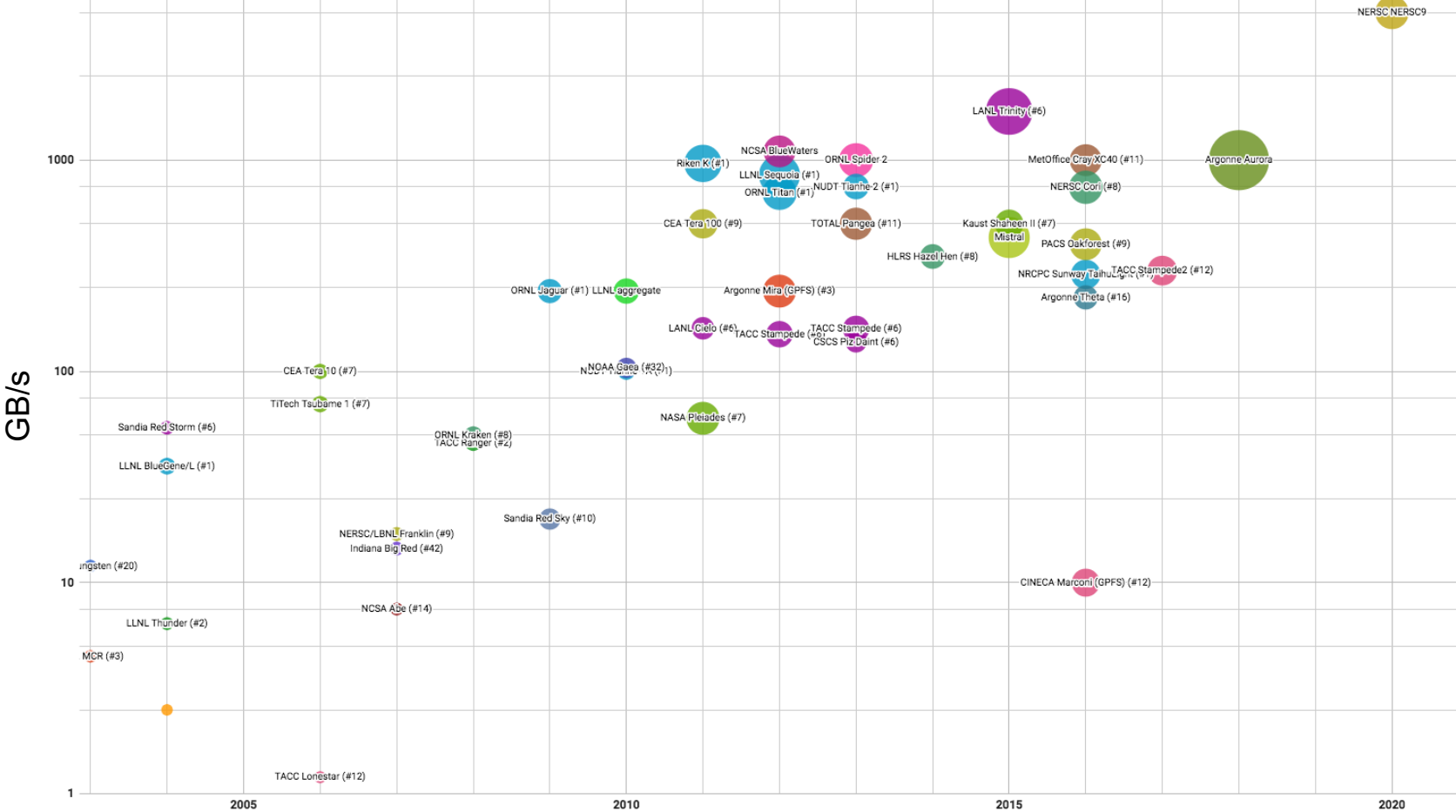


Lustre Feature Renaissance

Nathan Rutman
Lug 2018

14 Years of Scaling



Years of effort spent on performance and scaling

Not an exhaustive list

- Idiskfs scaling, ZFS
- Recovery: interop, VBR, AT, FSCK, Imperative Recovery
- CLIO, MDS rewrite, FIDs
- IO scaling: LRU, read cache, readahead, wide striping, multi-MB rpc, DoM
- MD scaling: statahead, DNE, MMR
- tons of diverse performance improvements
- bugs bugs bugs

Features too of course: mountconf, Kerberos, NRS, HSM, changelogs, pools, multirail



Capable infrastructures in place

- DNE – MD horizontal scaling
- Complex layouts – much more interesting data placement
- FLR – data redundancy inside Lustre

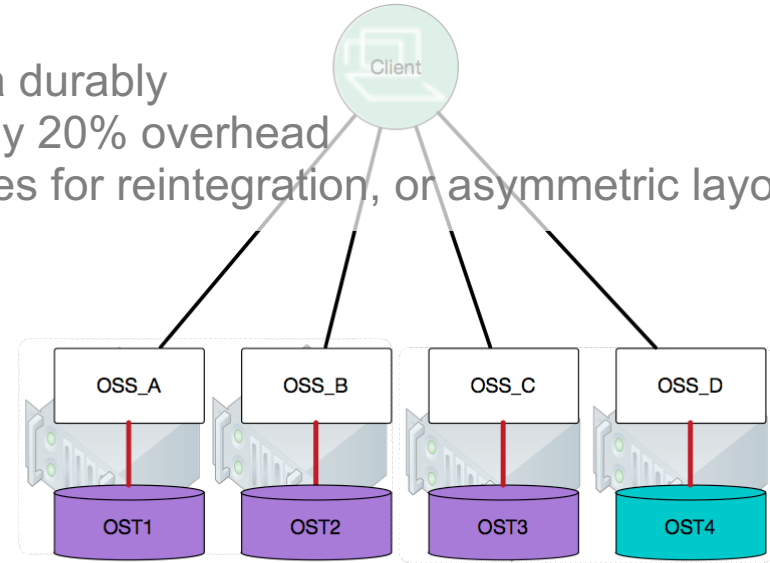
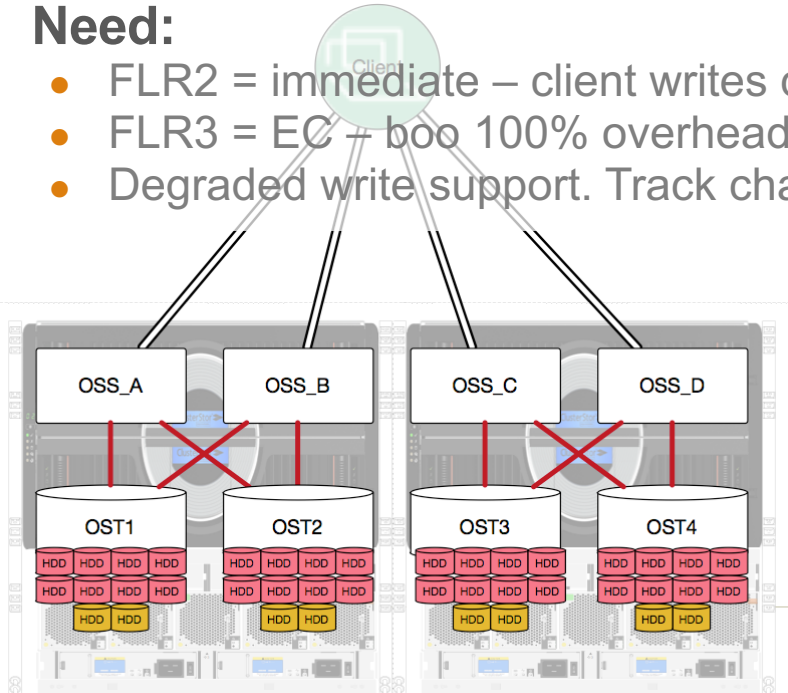
Time to reap some Feature rewards

- Let's look at some possible features
- These aren't even designs, just ideas

FLR – data redundancy *inside* Lustre

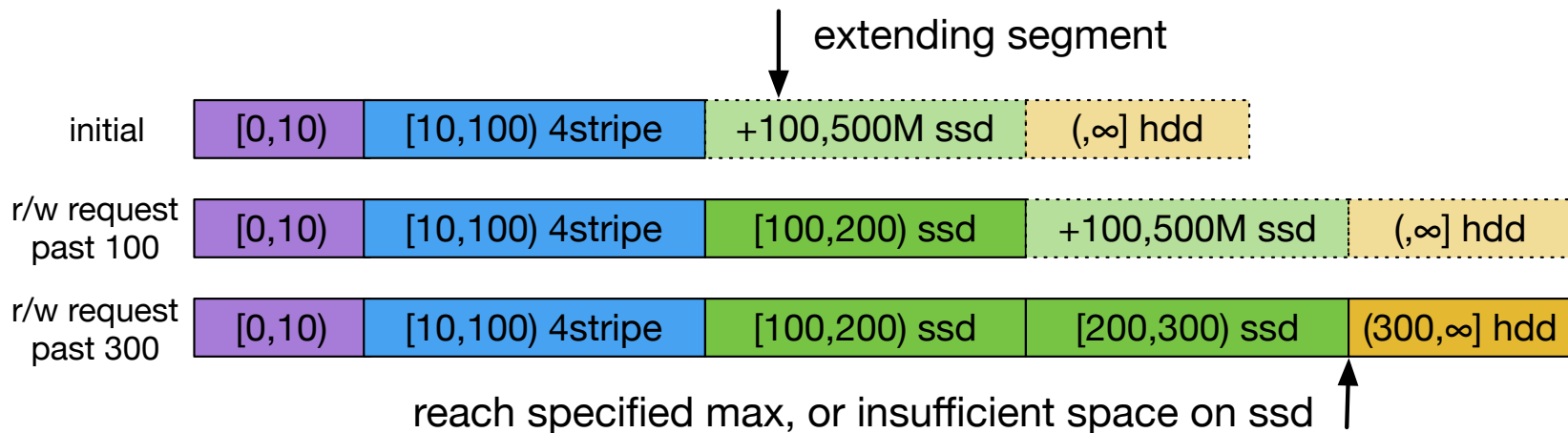
- One small (?) step for Layouts, one giant leap for Lustre systems design
- No longer need to rely on ~~Failover~~ for data access
- Dual-ported, dual-server, dual-path, dual-\$ - nope.
- Need:

- FLR2 = immediate – client writes data durably
- FLR3 = EC – boo 100% overhead, yay 20% overhead
- Degraded write support. Track changes for reintegration, or asymmetric layouts?



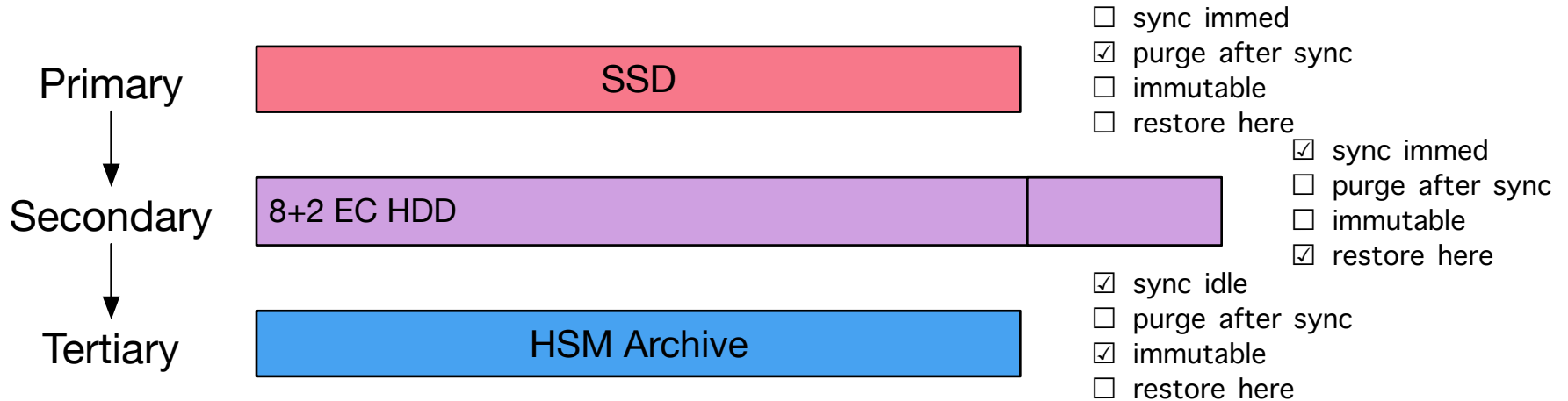
Spillover Space: death to ENOSPC

- Self-extending PFL (LU-10070, LU-10169)
- Some PFL segments are virtual, instantiated on demand
- Request in a virtual segment requires layout update
- MDS adds a new component on demand
- Can choose the new component striping based on dynamic conditions (e.g. free space)



ILM Layouts

- Layout implies an action: stale FLR copy = resync w/ lfs mirror
- And a timeframe: (immediate | eventual)
- Simple ILM policy *already* encoded into layout
- Add some flags to layout and/or policy ref
- Make HSM a true layout (LU-10606): stale HSM copy = resync w/ lfs hsm
- Use Coordinator and Copytool for all movement (LU-6081)



Asymmetric Layouts

- Reads go to R iff not in W
- Block bitmap on W tracks newly written data
- Client caching and DIO insure full-page writes
- W controls all locks, gives bitmap along with lock grants
- Clients access R or W directly, all under W's locks

Why?

- Write to flash, read from HDD
- Continue writing to new W if an OST fails (checkpoint) (or ENOSPC)
- EC degraded write case – point of EC is to remain usable in failures





Fast Find

- **Why do we copy Lustre MD into DB's or scan raw ldiskfs?**
- **Need to quickly find files that match certain criteria**
- **A great 'lfs find' could do the same thing, saving the tools effort**
 - Server side. RPC from client, returns filtered list
 - Logical combinations of filters
 - Unix-style piping: *lfs find /lustre -size +20M | lfs hsm archive*
- **Add new MDT indices to efficiently generate initial candidate lists**
 - LRU, file heat, mtime, size
 - dt_index_operations (eg IAM) provides generic indexing code
 - Update indices transactionally with MD updates



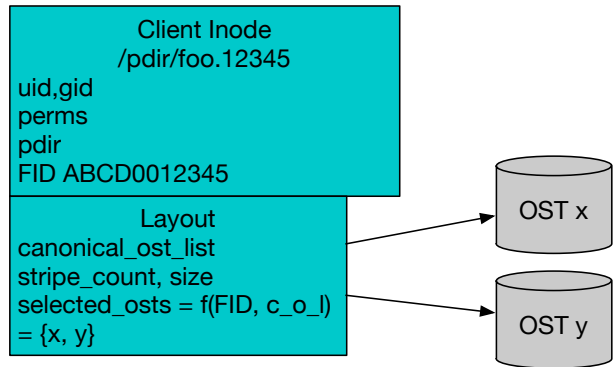
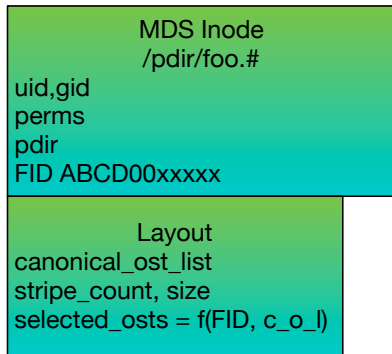
Rough SoM

- **FLR records file size on MDS; comes with sync**
- **DoM records file size on MDS**
- **Straightforward to get maximum size, if we don't care about evicted/failover case**
- **Rough size is fine for many purposes (e.g. policies)**
- **Record the quality of SoM, let users decide if usable**
- **Strict, Rough, Stale, Unknown (LU-9538)**
- **Don't return as POSIX size unless strict**

Clone Files: extreme create scaling



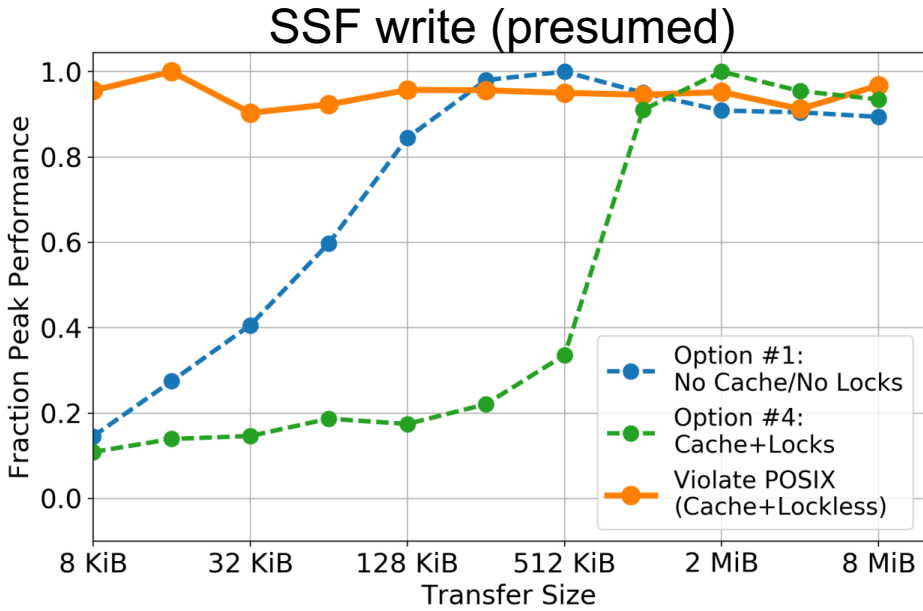
- File create: ask MDS to create, lock dir, create inode, assign objects
- Clone create: create “all” the files at once
- Single MDS inode, single namespace entry: foo.#
- FID is prefix+#
- Layout is f(FID)
- Shared MD (clones!) but different objects / data / sizes
- open(foo.4,O_CREAT) is now a client-local operation



Alternate Consistency Models



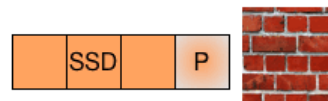
- **POSIX API vs POSIX consistency semantics**
- **Caching allow write coalesce, local latency**
- **But pay a penalty for locks**
- **Solutions in Lustre, but requires effort**
 - Lockless DIO, Grouplocks
- **Make it easier**
 - `ladvise`?
 - Persistent file tags?
 - Automatically change modes?



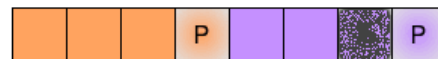
Glenn Lockwood
<https://www.nextplatform.com/2017/09/11/whats-bad-posix-io/>

All Together Now

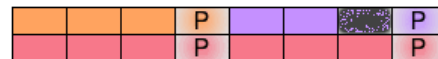
EC



Spillover




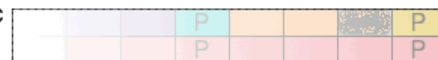
Asymmetric



ILM
sync on close

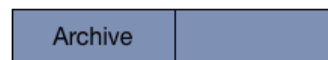


purge after sync



Fast Find
big (SoM), old files

HSM



Feature Vote?

1. FLR EC with degraded writes
- ~~2. Spillover Space~~ *cray*
3. ILM Layouts
4. Asymmetric Layouts
5. Fast Find
- ~~6. Rough SoM~~ *DDN*
7. Clone Files
8. Alternate Consistency Hints

Implementation Plan

1. Ignore Nathan's slideware
2. <insert smart developer here>



3. Implement!